Lecture 21: Simply-Typed Lambda Calculus

CS4400 Programming Languages

Syntax extensions:

• Note that abstractions need to specify the type of the bound variable – there is no way for the type-checker to guess it (at this stage)

```
data Expr = ...
| Lam Variable Type Expr
| App Expr Expr
data Type = ...
| TyArrow Type Type
```

TyArrow:

• The new *type constructor*, **TyArrow**, represents a function type:

TyArrow TyInt TyBool is the type a function that takes an integer (TyInt) and returns a boolean (TyBool). In Haskell (also in some other languages and in type theory), this is written Integer -> Bool

TyArrow (TyArrow TyInt TyBool) (TyArrow TyInt TyBool) corresponds to (Integer -> Bool) -> (Integer -> Bool), that is, the type of a function that takes a function from integers to booleans and returns a function from integers to booleans.

Due to currying, we normally understand this as a function that takes a function from integers to booleans, then an integer and returns a boolean. Note that this also means that the arrow -> is *right-associative* and the above Haskell type can be equivalently written as (Integer -> Bool) -> Integer -> Bool. Also note, that this is opposite of how application associates, which is to the left.

Note on associativity:

Function type – RIGHT: $t1 \rightarrow t2 \rightarrow t3 \rightarrow t4$ is the same as $t1 \rightarrow (t2 \rightarrow t3 \rightarrow t3)$ is the same as $t1 \rightarrow (t2 \rightarrow t3 \rightarrow t3)$ Function application – LEFT: f a b c is the same as (f a) b c is the same as ((f a) b) c

Rules

```
add x t1 tenv |- e : t2

tenv |- Lam x t1 e : TyArrow t1 t2

tenv |- e1 : TyArrow t2 t1 e2 : t2' t2 == t2'

tenv |- App e1 e2 : t1
```

The fixpoint operator:

• No fixpoint combinator (e.g., Y or Z) can be type-checked in STLC, so it has to be added as a primitive operation

data Expr = ... | Fix Expr